

CS 211: Linear Search, Binary Search, Sorting

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Week 12

Searching Linearly

- ▶ Return the first index of key in array arr
- ▶ Return -1 if not present
- ▶ No assumptions on ordering of arr
- ▶ What assumptions are needed about available methods?

For Ints

```
public static int linearSearch(int arr[], int key)
```

For Strings

```
public static int linearSearch(String arr[], String key)
```

For Foos

```
public static int linearSearch(Foo arr[], Foo key)
```

Linear Search for Primitives and Objects

```
public static
int linearSearch(int arr[],
                 int key)
{
    for(int i=0; i<arr.length; i++){
        if(arr[i] == key){
            return i;
        }
    }
    return -1;
}
```

```
public static
int linearSearch(Object arr[],
                 Object key)
{
    for(int i=0; i<arr.length; i++){
        if(arr[i].equals( key )){
            return i;
        }
    }
    return -1;
}
```

Questions

- ▶ Could you go faster with a sorted array?
- ▶ If so how?
- ▶ How much work is it to sort an array?

Binary Search

- ▶ A faster way to search **sorted** arrays
- ▶ Array **must** be sorted: linear search doesn't require that
- ▶ Like searching for a word in a dictionary: repeated halving of search space (look left or look right)
- ▶ **Let's write it:**

```
// Iterative
// Use binary search to locate an a given integer in an array
public static int binarySearch(int a[], int key)

// Easy-to-use recursive version which calls a helper
public static int binarySearchR(int a[], int key){
    return binarySearchR(a,key,0,a.length-1);
}

// Recursive
// Helper method which does work of searching, repeatedly
// halving search area by adjusting left/right
public static int binarySearchR(int a[], int key,
                                int left, int right)
```

Binary Search: Iterative

```
// Iterative: Use binary search to
// locate integer key in array a
public static
int binarySearch(int a[], int key){
    int left=0, right=a.length-1;
    int mid = 0;
    while(left <= right){
        mid = (left+right)/2;
        if(key == a[mid]){
            return mid;
        }else if(key < a[mid]){
            right = mid-1;
        }
        else{
            left = mid+1;
        }
    }
    return -1;
}
```

Demonstration

```
key: 6 -> index 2
key: 13 -> not found
key: 2 -> index 0
key: 25 -> not found
```

array a:

	0		1		2		3		4		5		6		7		8		9	
	-----+																			
	2		3		6		9		11		12		14		15		17		21	

Binary Search: Recursive

```
// Easy-to-use recursive version which calls a helper
public static int binarySearchR(int a[], int key){
    return binarySearchR(a,key,0,a.length-1);
}
// Helper method which does work of searching
public static int binarySearchR(int a[], int key,
                                int left, int right){
    if(left > right){
        return -1;
    }
    int mid = (left+right)/2;
    if(key == a[mid]){
        return mid;
    }else if(key < a[mid]){
        return
            binarySearchR(a,key,left,mid-1);
    }
    else{
        return
            binarySearchR(a,key,mid+1,right);
    }
}
```

Speed of Linear vs Binary Search

Specific Size

- ▶ Input array 1024 elements
- ▶ key is not present
- ▶ How many steps to determine key is not there?
 - ▶ For linear search?
 - ▶ For binary search?

General Size

- ▶ Input array N elements
- ▶ key is not present
- ▶ How many steps to determine key is not there?
 - ▶ For linear search?
 - ▶ For binary search?

Quick Review: Search

- ▶ Describe one way to search for data in an array; include any assumptions necessary for this process to work
- ▶ Describe a fundamentally different way
- ▶ Compare these two approaches according to their *worst case runtime complexity*

Sorting

For binary search to work, must have sorted input

- ▶ How do I get ints or Strings or anything else sorted?
- ▶ What do you know so far about sorting?
- ▶ What you should know as Computer Scientist is ...

How does sorting Actually Work?

- ▶ Tons of CPU time devoted to sorting
- ▶ Records in databases
- ▶ Rows/Columns in spreadsheets
- ▶ We will discuss **Selection Sort**, simple and inefficient
 - ▶ (BJP 13.3)
- ▶ May have time to discuss **Merge Sort** which is more efficient
 - ▶ (BJP 13.4)
- ▶ Get hungry for better sorting algorithms

Other courses

- ▶ CS 211: Start basic discussions
- ▶ CS 310: Framework for analysis (Big-O notation)
- ▶ CS 483: Analyze detailed algorithms

Selection Sort

- ▶ Dead simple sorting
- ▶ Repeatedly look for the minimum element in right part of array
- ▶ Swap min element with last element of left part of array

```
public static void selectionSort(int[] a) {
    for (int i = 0; i < a.length - 1; i++) {
        int smallest = i;
        for (int j = i + 1; j < a.length; j++) {
            if (a[j] < a[smallest]) {
                smallest = j;
            }
        }
        swap(a, i, smallest);
    }
}

public static void swap(int[] a, int i, int j) {
    int temp = a[i]; a[i] = a[j]; a[j] = temp;
}
```

An A/V Intensive Demonstration

- ▶ Selection Sort by Timo Bingmann, hosted on YouTube
- ▶ Also worth a watch: Gypsy Dance Demo of Selection Sort on YouTube

Demonstration of Selection Sort

```
a = [ 19 18 21 12 14 16 13 17 42 31]
      0  1  2  3  4  5  6  7  8  9
```

```
Outer i = 0, smallest = 0 (19)
```

```
Inner
```

```
j = 1, smallest = 1 (18)
```

```
j = 2, smallest = 1 (18)
```

```
j = 3, smallest = 3 (12)
```

```
j = 4, smallest = 3 (12)
```

```
...
```

```
j = 9, smallest = 3 (12)
```

```
swap(a, i, smallest)
```

```
a = [ 12 18 21 19 14 16 13 17 42 31]
      0  1  2  3  4  5  6  7  8  9
```

```
Outer i = 1, smallest = 1 (18)
```

```
Inner
```

```
j = 2, smallest = 1 (18)
```

```
j = 3, smallest = 1 (18)
```

```
j = 4, smallest = 4 (14)
```

```
j = 5, smallest = 4 (14)
```

```
...
```

```
j = 9, smallest = 6 (13)
```

```
swap(a, i, smallest)
```

```
a = [ 12 13 21 19 14 16 18 17 42 31]
      0  1  2  3  4  5  6  7  8  9
```

```
Outer i = 2, smallest = 2 (21) ...
```

```
public static
void selectionSort(int[] a) {
    for(int i=0; i<a.length-1; i++){
        int smallest = i;
        for(int j=i+1; j<a.length; j++){
            if (a[j] < a[smallest]) {
                smallest = j;
            }
        }
        swap(a, i, smallest);
    }
}
```

Show the next few outer iterations

Runtime of Selection sort

```
public static void selectionSort(int[] a) {
    for (int i = 0; i < a.length - 1; i++) {
        int smallest = i;
        for (int j = i + 1; j < a.length; j++) {
            if (a[j] < a[smallest]) {
                smallest = j;
            }
        }
        swap(a, i, smallest);
    }
}

public static void swap(int[] a, int i, int j) {
    int temp = a[i]; a[i] = a[j]; a[j] = temp;
}
```

- ▶ Each inner loop iteration does a constant amount of work: Fetch array elements, compare numbers, set variables
- ▶ Good approximation of runtime: total inner loop iterations
- ▶ **Total inner iterations** for array size 10? size 50? size N ?

Speed and Alternatives

Selection Sort is

- ▶ Easy to code, easy to understand
- ▶ Has terrible runtime performance
- ▶ $O(N^2)$ WORST case performance
- ▶ $O(N^2)$ BEST case performance
 - ▶ Pre-sorted array still takes $O(N^2)$ operations

Alternatives

- ▶ Many sorting algorithms exist
- ▶ [Visualization and Comparison of Sorting Algorithms by Viktor Bohush, hosted on YouTube](#)
- ▶ Good sorting algorithms have $O(N \log N)$ worst case performance
- ▶ What algorithm does `Arrays.sort()` and `Collections.sort()` use?

Asymptotic Speeds To Remember

Input Data

Array of size N

Search

- ▶ Linear search: $O(N)$
- ▶ Binary search: $O(\log N)$
- ▶ Hash tables: $O(1)$

Sort

- ▶ Selection sort: $O(N^2)$
- ▶ Insertion sort: $O(N^2)$
- ▶ Merge sort: $O(N \log N)$
- ▶ Quick sort: $O(N \log N)$
- ▶ Radix sort: $O(N)$ via cheating

Generalizing Binary Search

What would change in the binary search code if String were the operative type rather than int?

```
public static
int binarySearch(String a[],
                 String key)
```

```
public static
int binarySearch(int a[],
                 int key)
{
    int left=0, right=a.length-1;
    int mid = 0;
    while(left <= right){
        mid = (left+right)/2;
        if(key == a[mid]){
            return mid;
        }else if(key < a[mid]){
            right = mid-1;
        }
        else{
            left = mid+1;
        }
    }
    return -1;
}
```

Generalized Searching and Sorting

Search/Sort in the Java tied to two interfaces

`interface Comparable<T>` : Objects compare to each other

- ▶ Class has function `int compareTo(T y)`
`if(x.compareTo(y) < 0){...`
- ▶ `x.compareTo(y)`: Returns "x minus y"
 - ▶ Negative for x before y
 - ▶ 0 for equal
 - ▶ Positive for x after y
- ▶ Note: Not always -1/0/+1: **why not?**

`interface Comparator<T>` : Judge of two other objects

- ▶ Class has `compare(T x,T y)`: Returns "x minus y"
 - ▶ Neg/0/Pos numbers for ordering
- ```
Comparator<String> cmp = ...;
if(cmp.compare(x,y) < 0){...
```

## Quick Review: Sorting

### Sorting

- ▶ Describe how one sorting algorithm works
- ▶ Give its worst-case runtime complexity
- ▶ What is the runtime complexity of the best sorting algorithms?
- ▶ Given an example of one sorting algorithm that achieves this complexity
- ▶ What sorting algorithm does `Collections.sort(..)` use anyway?

### Comparing

- ▶ What's one way a java class can be made to work with functions like `Arrays.binarySearch(..)` and `Collections.sort(..)`?
- ▶ Are there any other ways?

# Generalized Binary Search

Adapt String version of binary search to

```
public static <T extends Comparable<T>>
int binarySearch(T a[], T key)
{
 ...
}
```

Works with any Comparable thing:  
String, Integer, Person

```
public static
int binarySearch(int a[],
 int key)
{
 int left=0, right=a.length-1;
 int mid = 0;
 while(left <= right){
 mid = (left+right)/2;
 if(key == a[mid]){
 return mid;
 }else if(key < a[mid]){
 right = mid-1;
 }
 else{
 left = mid+1;
 }
 }
 return -1;
}
```

## Visit your Local Library

Arrays and Collections have `sort()` methods for Comparable stuff and which take a Comparator

### From Arrays

```
// Uses compareTo of comparable
static void sort(Object[] a)
```

```
// Uses comparator
static <T> void sort(T[] a, Comparator<? super T> c)
```

### From Collections

```
// Uses compareTo of comparable
static <T extends Comparable<? super T>> void sort(List<T> list)
```

```
// Uses comparator
static <T> void sort(List<T> list, Comparator<? super T> c)
```

## Comparable Exercise

### Person

Write a class Person which implements the Comparable interface

```
class Person implements Comparable<Person>{
 public Person(String first, String last);
 public String toString();
 public int compareTo(Person other);
}
```

- ▶ Must have last method to satisfy Comparable interface
- ▶ Sort by **last name** then **first name**
- ▶ String already has a compareTo so this is easy

Try this with Generalized Binary Search

## Comparators and Sorting

- ▶ Sorting algs in Arrays and Collections are generalized
- ▶ Comparator provides a powerful way to sort in new ways
  - ▶ Just define `int compare(T x, T y)` and sort
- ▶ Example: Reverse numeric comparison:

```
import java.util.*;
class RevComp implements Comparator<Integer>{
 public int compare(Integer x, Integer y){
 return y-x;
 }
}
public class RevComparator{
 public static void main(String args[]){
 Integer a[] = {4, 5, 9, 2, 3, 1, 8};
 Comparator<Integer> cmp = new RevComp();
 Arrays.sort(a, cmp);
 System.out.println(Arrays.toString(a));
 // [9, 8, 5, 4, 3, 2, 1]
 }
}
```

## Exercise: Odds First

Define a Comparator on Integers which sorts odds before evens with odd numbers in order within the first section and even numbers in order within the second section.

```
public static void main(String args[]){
 Comparator<Integer> cmp = new OddsThenEvens();
 Integer a[];

 a = new Integer[]{4, 5, 9, 2, 3, 1, 8};
 System.out.println(Arrays.toString(a));
 Arrays.sort(a,cmp);
 System.out.printf("%s\n\n",Arrays.toString(a));
 // [1, 3, 5, 9, 2, 4, 8]

 a = new Integer[]{2048, 1024, 5096, 128, 9999};
 System.out.println(Arrays.toString(a));
 Arrays.sort(a,cmp);
 System.out.printf("%s\n\n",Arrays.toString(a));
 // [9999, 128, 1024, 2048, 5096]
}
```



## Optional: Merge Sort

- ▶ Fast sorting algorithm:  $O(N \log N)$
- ▶ Exploits recursion
- ▶ Principles are simple but implementation is non-trivial
- ▶ Variants have different properties: out-of place vs in-place
- ▶ Good culmination of CS 211 last two weeks of material

## Exercise: Merge

```
public static void merge(int[] result, int[] a, int[] b)
```

- ▶ a and b are sorted arrays, may not be same size
- ▶ Copy elements from a and b into result so that result is sorted
- ▶ Assume: `result.length = a.length + b.length`
- ▶ Target Runtime Complexity:  $O(N)$  where  $N$  is `result.length`

### Example

```
int [] a = {1, 6, 7, 9}
int [] b = {0, 2, 3, 4, 8}
int [] result = new int[a.length+b.length];

merge(result, a, b);

print(result)
// [0, 1, 2, 3, 4, 6, 7, 8, 9]
```

# Merge Answers

## Textbook

```
public static
void merge(int[] result,
 int[] a, int[] b) {
 int i1 = 0;
 int i2 = 0;
 for(int i = 0; i < result.length; i++){
 if(i2 >= b.length ||
 (i1 < a.length && a[i1] <= b[i2])){
 result[i] = a[i1];
 i1++;
 }
 else {
 result[i] = b[i2];
 i2++;
 }
 }
}
```

## CK preferred

```
public static
void merge(int[] result,
 int[] a, int[] b) {
 int ai=0, bi=0;
 for(int ri=0; ri<result.length; ri++){
 if(ai >= a.length){
 result[ri] = b[bi];
 bi++;
 }
 else if(bi >= b.length){
 result[ri] = a[ai];
 ai++;
 }
 else if(a[ai]<=b[bi]){
 result[ri] = a[ai];
 ai++;
 }
 else{
 result[ri] = b[bi];
 bi++;
 }
 }
}
```

## Recursive Application

```
public static void mergeSort(int[] a) {
 if (a.length <= 1) {
 return;
 }
 int[] left = Arrays.copyOfRange(a, 0, a.length/2);
 int[] right = Arrays.copyOfRange(a, (a.length/2), a.length);

 mergeSort(left);
 mergeSort(right);

 merge(a, left, right);
}
```

- ▶ An array of 1 element is sorted
- ▶ If bigger, chop array into two halves
- ▶ Recursively sort left and right halves
- ▶ Merge the sorted results into the original array

## Demonstrate

```
public static void mergeSort(int[] a) {
 if (a.length <= 1) {
 return;
 }
 int[] left = Arrays.copyOfRange(a, 0, a.length/2);
 int[] right = Arrays.copyOfRange(a, (a.length/2), a.length);

 mergeSort(left);
 mergeSort(right);

 merge(a, left, right);
}
```

## Show Execution for

```
int [] a = {14, 32, 67, 76, 23, 41, 58, 85};
mergeSort(list);
```

`mergeSortDebug()` in `CKMergeSort.java` is useful to see this

# Computational Complexity

- ▶ What is the complexity of `mergeSort()`
- ▶ Could you **prove** it?
- ▶ What's one huge disadvantage of this version of merge sort?

## In-Place Merge Sort

- ▶ Making copies of arrays takes time and memory
- ▶ Current version is an out-of-place merge sort
- ▶ For an array of size 1,000,000, need `left` and `right` arrays which total another 1,000,000 units of memory - BAD!
- ▶ Prefer an in-place version to save memory
- ▶ Cost: **Implementation complexity**
- ▶ Examine: [In-place merge sort ported from C++ STL](#)