

**Solutions**

1. [50 points]

Prove the following equivalence by substitution, i.e., use known logical equivalences to show that  $\neg(p \vee \neg q) \vee (\neg p \wedge \neg q)$  is equivalent to  $\neg p$ . *You must start from the statement  $\neg(p \vee \neg q) \vee (\neg p \wedge \neg q)$ .* Justify each step with the name of the corresponding logical equivalence being used. For full credit, do *not* skip steps.

$$\neg(p \vee \neg q) \vee (\neg p \wedge \neg q) \equiv \neg p$$

$\neg(p \vee \neg q) \vee (\neg p \wedge \neg q) \equiv (\neg p \wedge q) \vee (\neg p \wedge \neg q)$	De Morgan's law
$\equiv \neg p \wedge (q \vee \neg q)$	distributive law
$\equiv \neg p \wedge \text{TRUE}$	excluded middle law
$\equiv \neg p$	property of the constant TRUE

2. [50 points]

Using rules of inference with no substitution, prove that *modus ponens* is a valid rule. Use the notation introduced in class, and state, for each line, the rule of inference that justifies it. Make sure you do *not* to use *modus ponens* itself in your proof!

*Modus ponens:*                      **Write below the equivalent expression you need to prove:**

$p \rightarrow q$	
$p$	
_____	
$q$	

**Equivalent expression:**

$$((p \rightarrow q) \wedge p) \rightarrow q$$

Proof:

$[(p \rightarrow q) \wedge p]$	Assumption
$p \rightarrow q$	$\wedge$ elimination
$p$	$\wedge$ elimination
$[\neg q]$	Assumption
$\neg p$	Modus tollens
FALSE	Contradiction
$q$	Reduction to absurdity
$((p \rightarrow q) \wedge p) \rightarrow q$	$\rightarrow$ introduction