# Lecture: Analysis of Algorithms (CS583-004) 

Amarda Shehu

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(1) Graphs

- Definition of a Graph
- Omnipresence of Graphs
(2) Graph Representations
- Adjacency List Representation
- Adjacency Matrix Representation
- Alternative Graph Representations
(3) Solving Problems with Graph Algorithms


## What is a Graph?

Graph $G=(V, E)$

- $V$ : set of vertices
- $E$ : set of edges consisting of pairs of vertices from $V$



## First Graph Problem

## Seven Bridges of Koenigsberg [1736]:

Find a route that crosses each bridge exactly once.
Posed by Leonard Euler [1707-1783].
 modified from wikipedia

What is the minimum number of bridges that need to be added so that there exists a route that crosses each bridge exactly once?

## Road Networks as Graphs



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## Airline Routes as Graphs



Figure: http://www.airlineroutemaps.com/

## Social Networks as Graphs



Figure: http://hbr.idnet.net/images/

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## The Internet as a Graph



Figure: Credit: Matt Britt

## Visualization of the various

 routes through a portion of the Internet.

Figure: Credit: Young Hyun, CAIDA

## Websites as Graphs



Figure: http://www.google.com

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blue: for links
red: for tables
green: for the DIV tag
violet: for images
            : for forms
orange: for linebreaks and blockquotes
black: the HTML tag, the root node
gray: all other tags
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Figure: Credit: Marcel Salathe http://www.aharef.info


Figure: http://www.cs.gmu.edu


Figure: http://www.apple.com

## Biological Networks as Graphs



Figure: Adapted from A. Barabasi, University of Notre Dame

## Applications of Graphs

- Compilers
- Databases
- Neural Networks
- Machine Learning
- Artificial Intelligence
- Robotics
- Computational Biology
- ...


## Formal Definition of a Graph

- A graph $G=(V, E)$ is a pair consisting of:
- a set $V$ of vertices (or nodes)
- a set $E \subseteq V \times V$ of edges (or arcs)
- edge $e_{i} \in E$ is a pair $(u, v)$ connecting vertices $u$ and $v$

A graph $G=(V, E)$ is:

- directed (referred to as a digraph) if $E$ is a set of ordered pairs of vertices. The edges here are often referred to as directed edges or arrows.
- undirected if $E$ is a set of unordered pairs of vertices.
- weighted if there are weights associated with the edges.


## Illustrations of Types of Graphs



Figure: undirected graph


Figure: multigraph


Figure: directed graph


Figure: weighted graph

## General Definition of a Graph

## In a graph $G=(V, E)$ :

- E may be a set of unorderered pairs of vertices not necessarily distinct. More than one edge can connect two vertices.
- An edge in $E$ may connect more than two vertices.
- These graphs are referred to as multigraphs or pseudo-graphs.


## Focusing on Simple Graphs

## Simple Graphs

- A simple graph, or a strict graph, is an unweighted, undirected graph containing no loops or multiple edges
- Given that $E \subseteq V \times V,|E| \in O\left(|V|^{2}\right)$.
- If a graph is connected, $|E| \geq|V|-1$
- Combining the two, show that $\lg (|E|) \in \theta(\lg (|V|))$


## More Definitions, Conventions, Nomenclature

- A subgraph $H$ of $G=(V, E)$ is $H=\left(V_{1}, E_{1}\right)$ where $V_{1} \subseteq V$ and $E_{1} \subseteq E$, where $\forall e=(k, j) \in E_{1}, k, j \in V_{1}$.
- A path is a sequence of vertices, where each pair of successive vertices is connected by an edge.
- The length of the path is the number of edges in the path.
- A simple path contains unique vertices.
- A cycle is a simple path with the same first and last vertex.
- Two vertices are adjacent if they are connected by an edge.
- The neighbors of a vertex are all the vertices adjacent to it.
- The degree of a vertex is the number of its neighbors.
- A graph is connected if $\exists$ a path between every pair of vertices.
- A tree is a connected graph with no cycles.


## Graph Representations

- A graph can be represented as an adjacency list.
- A graph can be represented as an adjacency matrix.


## Adjacency List Representation



## Basic Graph Functionality

| Function | Adjacency List |  |
| :---: | :---: | :---: |
| find ( $v$ ) | $O(\|V\|)$ |  |
| hasVertex( $v$ ) | $O($ find $(\mathrm{v})$ ) |  |
| hasEdge( $v_{i}, v_{j}$ ) | $O\left(\right.$ find $\left.\left(\mathrm{v}_{\mathrm{i}}\right)+\operatorname{deg}\left(\mathrm{v}_{\mathrm{i}}\right)\right)$ |  |
| insertVertex( $v$ ) | $O(1)$ | $\|\operatorname{elist}[\mathrm{v}]\|=\operatorname{degree}(v) .$ |
| insertEdge( $v_{i}, v_{j}$ ) | $O\left(\right.$ find $\left.\left(\mathrm{v}_{\mathrm{i}}\right)\right)$ |  |
| removeVertex ( $v$ ) | $O(\|V\|+\|E\|)$ | In digraphs: |
| removeEdge( $v_{i}, v_{j}$ ) | $O\left(\right.$ find $\left.\left(\mathrm{v}_{\mathrm{i}}\right)+\operatorname{deg}\left(\mathrm{v}_{\mathrm{i}}\right)\right)$ | $\|\operatorname{list}[\mathrm{v}]\|=$ out-degree(v). |
| outEdges(v) | $O($ find $(\mathrm{v})+\operatorname{deg}(\mathrm{v}))$ |  |
| inEdges( $v$ ) | $O(\|V\|+\|E\|)$ |  |
| overall memory | $O(\|V\|+\|E\|)$ |  |

Handshaking Lemma: $\sum_{v \in V}|\operatorname{elist}(\mathrm{v})|=2|\mathrm{E}|$ for undirected graphs. $O(|V|+|E|)$ storage $\Rightarrow$ sparse representation.

## More on Implementation of Adjacency-list Representation

The adjacency list of a vertex can be implemented as a linked list The list of vertices themselves can be implemented using:

- A linked list
- A binary search tree
- A hash table

In a standard implementation, each edge list has two fields, a data field and a pointer:

- The data field contains adjacent vertex name and edge information
- The pointer points to next adjacent vertex


## Adjacency Matrix Representation


$M[i][j]=1 \quad$ iff $\quad\left(v_{i}, v_{j}\right) \in E$

| M | $V_{0}$ | $V_{1}$ | $V_{2}$ | $V_{3}$ | $V_{4}$ |  | M | $V_{0}$ | $V_{1}$ | $V_{2}$ | $V_{3}$ | $V_{4}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $V_{0}$ | 0 | 1 | 0 | 1 | 0 |  | $V_{0}$ | 0 | 1 | 1 | 0 | 1 |
| $v_{1}$ |  | 0 | 1 | 0 | 1 | bool **M; | $V_{1}$ | 0 | 1 | 0 | 0 | 0 |
| $V_{2}$ |  |  | 1 | 0 | 0 |  | $V_{2}$ | 0 | 1 | 0 | 1 | 0 |
| $V_{3}$ |  |  |  | 0 | 0 | using namespace std; | $v_{3}$ | 0 | 0 | 0 | 0 | 1 |
| $V_{4}$ |  |  |  |  | 1 | $\begin{aligned} & \text { vector < vector<bool\gg } \\ & \mathrm{M} \text {; } \end{aligned}$ | $v_{4}$ | 1 | 0 | 1 | 1 | 0 |

## Basic Graph Functionality

| Function | Adjacency Matrix |
| :--- | :--- |
| find $(v)$ | $O(1)$ |
| hasVertex $(v)$ | $O(1)$ |
| hasEdge( $\left(v_{i}, v_{j}\right)$ | $O(1)$ |
| insertVertex $(v)$ | $O\left(\|V\|^{2}\right)$ |
| insertEdge $\left(v_{i}, v_{j}\right)$ | $O(1)$ |
| removeVertex $(v)$ | $O\left(\|V\|^{2}\right)$ |
| removeddge $\left(v_{i}, v_{j}\right)$ | $O(1)$ |
| outEdges $(v)$ | $O(\|V\|)$ |
| inEdges $(v)$ | $O(\|V\|)$ |
| overall memory | $O\left(\|V\|^{2}\right)$ |

$O\left(|V|^{2}\right)$ storage $\Rightarrow$ dense representation.

## Comparing The Two Representations

| Function | Adjacency List | Adjacency Matrix |
| :--- | :--- | :--- |
| find $(v)$ | $O(\|V\|)$ | $O(1)$ |
| hasVertex $(v)$ | $O($ find $(\mathrm{v}))$ | $O(1)$ |
| hasEdge $\left(v_{i}, v_{j}\right)$ | $O\left(\right.$ find $\left.\left(\mathrm{v}_{\mathrm{i}}\right)+\operatorname{deg}\left(\mathrm{v}_{\mathrm{i}}\right)\right)$ | $O(1)$ |
| insertVertex $(v)$ | $O(1)$ | $O\left(\|V\|^{2}\right)$ |
| insertEdge $\left(v_{i}, v_{j}\right)$ | $O\left(\right.$ find $\left.\left(\mathrm{v}_{\mathrm{i}}\right)\right)$ | $O(1)$ |
| removeVertex $(v)$ | $O(\|V\|+\|E\|)$ | $O\left(\|V\|^{2}\right)$ |
| removeEdge $\left(v_{i}, v_{j}\right)$ | $O\left(\operatorname{find}\left(\mathrm{v}_{\mathrm{i}}\right)+\operatorname{deg}\left(\mathrm{v}_{\mathrm{i}}\right)\right)$ | $O(1)$ |
| outEdges $(v)$ | $O($ find $(\mathrm{v})+\operatorname{deg}(\mathrm{v}))$ | $O(\|V\|)$ |
| inEdges $(v)$ | $O(\|V\|+\|E\|)$ | $O(\|V\|)$ |
| overall memory | $O(\|V\|+\|E\|)$ | $O\left(\|V\|^{2}\right)$ |

## Alternative Graph Representations

## HashMap

Fast to query [hasVertex, hasEdge]
Fast to scan [outEdges]
$O(|V|)$
Fast to insert [insertVertex, insertEdge]
$O$ (1)
Fast to remove [removeEdge]
$O$ (1)

## Graph Representation: Hash Map

- Vertex set as a hash map
- key: vertex
- data: outgoing edges
- Outgoing edges of each vertex as a hash set



## Comparing The Three Representations

| Function | Adj. List | Adj. Matrix | Hash Map |
| :--- | :--- | :--- | :--- |
| find $(v)$ | $O(\|V\|)$ | $O(1)$ | $O(1)$ |
| hasVertex $(v)$ | $O(\|V\|)$ | $O(1)$ | $O(1)$ |
| hasEdge( $\left(v_{i}, v_{j}\right)$ | $O\left(\|V\|+\operatorname{deg}\left(v_{i}\right)\right)$ | $O(1)$ | $O(1)$ |
| insertVertex $(v)$ | $O(1)$ | $O\left(\|V\|^{2}\right)$ | $O(1)$ |
| insertEdge $\left(v_{i}, v_{j}\right)$ | $O(\|V\|)$ | $O(1)$ | $O(1)$ |
| removeVertex $(v)$ | $O(\|V\|+\|E\|)$ | $O\left(\|V\|^{2}\right)$ | $O(\|V\|)$ |
| removeEdge $\left(v_{i}, v_{j}\right)$ | $O\left(\|V\|+\operatorname{deg}\left(v_{i}\right)\right)$ | $O(1)$ | $O(1)$ |
| outEdges $(v)$ | $O(\|V\|+\operatorname{deg}(v))$ | $O(\|V\|)$ | $O(\operatorname{deg}(v))$ |
| inEdges $(v)$ | $O(\|V\|+\|E\|)$ | $O(\|V\|)$ | $O(\|V\|)$ |
| overall memory | $O(\|V\|+\|E\|)$ | $O\left(\|V\|^{2}\right)$ | linear-quadratic |

## Graph modeling: Problem Solving with Graph Algorithms

- Identify the vertices and the edges in your problem formulation
- Identify the objective of the problem
- State this objective in graph terms
- Implementation:
- Construct the graph from the input instance
- Run the suitable graph algorithm on the graph
- Convert the output into a suitable/required format

