### Lecture: Analysis of Algorithms (CS583 - 004)

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Amarda Shehu Lecture: Analysis of Algorithms (CS583 - 004)

#### 1 Graphs

- Definition of a Graph
- Omnipresence of Graphs

#### 2 Graph Representations

- Adjacency List Representation
- Adjacency Matrix Representation
- Alternative Graph Representations

#### 3 Solving Problems with Graph Algorithms

Definition of a Graph Omnipresence of Graphs

### What is a Graph?

#### Graph G = (V, E)

- V : set of vertices
- E : set of edges consisting of pairs of vertices from V



Definition of a Graph Omnipresence of Graphs

## First Graph Problem

#### Seven Bridges of Koenigsberg [1736]:

Find a route that crosses each bridge exactly once. Posed by Leonard Euler [1707 - 1783].





modified from wikipedia

What is the minimum number of bridges that need to be added so that there exists a route that crosses each bridge exactly once?

 $V_1$ 

Definition of a Graph Omnipresence of Graphs

#### Road Networks as Graphs



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#### Airline Routes as Graphs



Figure: http://www.airlineroutemaps.com/

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### Social Networks as Graphs



Figure: http://hbr.idnet.net/images/

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Definition of a Graph Omnipresence of Graphs

#### The Internet as a Graph



Figure: Credit: Matt Britt

Visualization of the various routes through a portion of the Internet.



Figure: Credit: Young Hyun, CAIDA

Definition of a Graph Omnipresence of Graphs

### Websites as Graphs



Figure: http://www.google.com

blue: for links red: for tables green: for the DIV tag violet: for images prime: for forms orange: for linebreaks and blockquotes black: the HTML tag, the root node gray: all other tags

Figure: Credit: Marcel Salathe http://www.aharef.info



Definition of a Graph Omnipresence of Graphs

### **Biological Networks as Graphs**



Figure: Adapted from A. Barabasi, University of Notre Dame

Definition of a Graph Omnipresence of Graphs

# Applications of Graphs

- Compilers
- Databases
- Neural Networks
- Machine Learning
- Artificial Intelligence
- Robotics
- Computational Biology
- ...

Definition of a Graph Omnipresence of Graphs

# Formal Definition of a Graph

- A graph G = (V, E) is a pair consisting of:
- a set V of vertices (or nodes)
- a set  $E \subseteq V \times V$  of edges (or arcs)
  - edge  $e_i \in E$  is a pair (u, v) connecting vertices u and v

#### A graph G = (V, E) is:

- directed (referred to as a digraph) if *E* is a set of ordered pairs of vertices. The edges here are often referred to as directed edges or arrows.
- undirected if E is a set of unordered pairs of vertices.
- weighted if there are weights associated with the edges.

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## Illustrations of Types of Graphs







Figure: directed graph



Figure: weighted graph

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## General Definition of a Graph

#### In a graph G = (V, E):

- *E* may be a set of unorderered pairs of vertices not necessarily distinct. More than one edge can connect two vertices.
- An edge in *E* may connect more than two vertices.
- These graphs are referred to as multigraphs or pseudo-graphs.

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## Focusing on Simple Graphs

#### Simple Graphs

- A simple graph, or a strict graph, is an unweighted, undirected graph containing no loops or multiple edges
- Given that  $E \subseteq V \times V$ ,  $|E| \in O(|V|^2)$ .
- If a graph is connected,  $|E| \geq |V| 1$
- Combining the two, show that  $lg(|E|) \in \theta(lg(|V|))$

Definition of a Graph Omnipresence of Graphs

## More Definitions, Conventions, Nomenclature

- A subgraph H of G = (V, E) is  $H = (V_1, E_1)$  where  $V_1 \subseteq V$ and  $E_1 \subseteq E$ , where  $\forall e = (k, j) \in E_1$ ,  $k, j \in V_1$ .
- A **path** is a sequence of vertices, where each pair of successive vertices is connected by an edge.
- The length of the path is the number of edges in the path.
- A simple path contains unique vertices.
- A cycle is a simple path with the same first and last vertex.
- Two vertices are **adjacent** if they are connected by an edge.
- The neighbors of a vertex are all the vertices adjacent to it.
- The **degree** of a vertex is the number of its neighbors.
- A graph is **connected** if ∃ a path between every pair of vertices.
- A tree is a connected graph with no cycles.

Adjacency List Representation Adjacency Matrix Representation Alternative Graph Representations

### Graph Representations

- A graph can be represented as an adjacency list.
- A graph can be represented as an adjacency matrix.

Adjacency List Representation Adjacency Matrix Representation Alternative Graph Representations

## Adjacency List Representation



Adjacency List Representation Adjacency Matrix Representation Alternative Graph Representations

## **Basic Graph Functionality**

Function	Adjacency List	
find(v)	O( V )	
hasVertex(v)	$O(\mathrm{find}(\mathrm{v}))$	
$hasEdge(v_i, v_j)$	$O(\operatorname{find}(v_i) + \operatorname{deg}(v_i))$	In undirected graphs:
insertVertex(v)	O(1)	$\left  \text{olign}(y) \right  = \text{degree}(y)$
$insertEdge(v_i, v_j)$	$O(\mathrm{find}(\mathrm{v_i}))$	$ \operatorname{enst}[v]  = \operatorname{degree}(v)$ .
removeVertex( $v$ )	O( V  +  E )	In digraphs:
removeEdge $(v_i, v_j)$	$O(\operatorname{find}(v_i) + \operatorname{deg}(v_i))$	elist[v]  = out-degree(v).
outEdges(v)	$O(\operatorname{find}(v) + \operatorname{deg}(v))$	
inEdges(v)	O( V + E )	
overall memory	O( V  +  E )	

Handshaking Lemma:  $\sum_{v \in V} |elist(v)| = 2|E|$  for undirected graphs. O(|V| + |E|) storage  $\Rightarrow$  sparse representation.

Adjacency List Representation Adjacency Matrix Representation Alternative Graph Representations

More on Implementation of Adjacency-list Representation

The adjacency list of a vertex can be implemented as a linked list The list of vertices themselves can be implemented using:

- A linked list
- A binary search tree
- A hash table

In a standard implementation, each edge list has two fields, a data field and a pointer:

- The data field contains adjacent vertex name and edge information
- The pointer points to next adjacent vertex

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## Adjacency Matrix Representation



Outline of Today's Class Graphs Solving Problems with Graph Algorithms

Adjacency List Representation Adjacency Matrix Representation Alternative Graph Representations

## Basic Graph Functionality

Adjacency Matrix
<i>O</i> (1)
O(1)
O(1)
$O( V ^2)$
O(1)
$O( V ^2)$
O(1)
O( V )
O( V )
$O( V ^2)$

 $O(|V|^2)$  storage  $\Rightarrow$  **dense** representation.

Adjacency List Representation Adjacency Matrix Representation Alternative Graph Representations

## Comparing The Two Representations

Function	Adjacency List	Adjacency Matrix
find(v)	O( V )	<i>O</i> (1)
hasVertex(v)	$O(\mathrm{find}(\mathrm{v}))$	O(1)
$hasEdge(v_i, v_j)$	$O(\operatorname{find}(v_i) + \operatorname{deg}(v_i))$	O(1)
insertVertex(v)	<i>O</i> (1)	$O( V ^2)$
$insertEdge(v_i, v_j)$	$O(\text{find}(v_i))$	O(1)
removeVertex( $v$ )	O( V  +  E )	$O( V ^2)$
removeEdge $(v_i, v_j)$	$O(\operatorname{find}(v_i) + \operatorname{deg}(v_i))$	O(1)
outEdges(v)	$O(\operatorname{find}(v) + \operatorname{deg}(v))$	O( V )
inEdges(v)	O( V + E )	O( V )
overall memory	O( V + E )	$O( V ^2)$

Adjacency List Representation Adjacency Matrix Representation Alternative Graph Representations

### Alternative Graph Representations

HashMap			
Fast to query	[hasVertex, hasEdge]	O(1)	
Fast to scan	[outEdges]	O( V )	
Fast to insert	[insertVertex, insertEdge]	O(1)	
Fast to remove	[removeEdge]	<i>O</i> (1)	

Adjacency List Representation Adjacency Matrix Representation Alternative Graph Representations

## Graph Representation: Hash Map

- Vertex set as a hash map
  - key: vertex
  - data: outgoing edges
- Outgoing edges of each vertex as a hash set



using namespace std\_ext;  
hash\_map  
vertex outgoing edges  
$$v_0$$
 hash\_set:  $v_1$ ,  $v_2$ ,  $v_4$   
hash\_set:  $v_1$   
hash\_set:  $v_1$   
hash\_set:  $v_2$   
hash\_set:  $v_4$   
hash\_set:  $v_4$   
hash\_set:  $v_4$   
hash\_set:  $v_7$ 

Adjacency List Representation Adjacency Matrix Representation Alternative Graph Representations

## Comparing The Three Representations

Function	Adj. List	Adj. Matrix	Hash Map
find(v)	O( V )	O(1)	<i>O</i> (1)
hasVertex(v)	O( V )	O(1)	O(1)
$hasEdge(v_i, v_j)$	$O( V  + \deg(v_i))$	O(1)	O(1)
insertVertex(v)	<i>O</i> (1)	$O( V ^2)$	O(1)
$insertEdge(v_i, v_j)$	O( V )	O(1)	O(1)
removeVertex(v)	O( V + E )	$O( V ^2)$	O( V )
$removeEdge(v_i, v_j)$	$O( V  + \deg(v_i))$	O(1)	O(1)
outEdges(v)	$O( V  + \deg(v))$	O( V )	$O(\deg(v))$
inEdges(v)	O( V + E )	O( V )	O( V )
overall memory	O( V + E )	$O( V ^2)$	linear-quadratic

# Graph modeling: Problem Solving with Graph Algorithms

- Identify the vertices and the edges in your problem formulation
- Identify the objective of the problem
- State this objective in graph terms
- Implementation:
  - Construct the graph from the input instance
  - Run the suitable graph algorithm on the graph
  - Convert the output into a suitable/required format